

Super-Time-Stepping

a simple way of speeding up the explicit scheme for parabolic problems:

Think of the PDE as $\frac{d}{dt} U(t) + A U(t) = 0$, $U(0) = u_0$

Think of the discrete scheme as $U^{n+1} = U^n - \Delta t_n A U^n = (I - \Delta t A) U^n$

$$U^0 = u_0$$

A is $M \times M$ matrix

symmetric, positive-definite
(for parabolic PDE)

Stability condition: $\rho(I - \Delta t A) < 1$, $\rho(\cdot)$ = spectral radius = $\max |\lambda_i|$

$$\Rightarrow \Delta t < \Delta t_{\text{expl}} := \frac{2}{\lambda_{\max}}, \quad \lambda_{\max} = \text{largest eigenvalue of } A$$

(CFL condition) $\left(= \frac{4D}{\Delta x^2} \text{ for } u_t = D u_{xx} \right)$

Idea: Do NOT require stability at each timestep Δt ,
only require it at the end of every N steps, for some chosen N
(1, 5, 10, 20, 100, ...)

So, consider a superstep ΔT consisting of N timesteps $\tau_1, \tau_2, \dots, \tau_N$
and want to choose $\{\tau_i\}$ to ensure stability over $\Delta T = \sum \tau_i$
and maximize the duration ΔT of the superstep.

So, want $\rho\left(\prod_{i=1}^N (I - \tau_i A)\right) < 1$ and $\sum_{i=1}^N \tau_i$ as large as possible.

Miracle: This optimization problem has explicit solution, thanks to the remarkable optimality properties of Chebychev polynomials !!!

$$\tau_i = \Delta t_{\text{expl}} \cdot \frac{1}{(-1+\nu) \cdot \cos\left(\frac{2i-1}{N} \frac{\pi}{2}\right) + 1+\nu}, \quad i=1, 2, \dots, N.$$

where $0 < \nu < \lambda_{\min}/\lambda_{\max}$

$$\text{Then } \Delta T \xrightarrow[\nu \rightarrow 0]{} \frac{N^2 \cdot \Delta t_{\text{expl}}}{N \cdot (N \cdot \Delta t_{\text{expl}})}$$

so N substeps of a superstep cover a time interval N times longer than N explicit steps (each of Δt_{expl}) !!!

That's where speedup comes from!

Implementation: Figure out Δt_{expl} from CFL condition (Positive Coeff. Rule).
Pick N, ν (say $N=5, \nu=0.001$).

Instead of executing steps of length Δt_{expl} ,
execute N steps of lengths $\tau_1, \tau_2, \dots, \tau_N$
(no outputting until a superstep is done)

superstep loop:

for $i=1:N$

$\tau_i = \dots$ (better: precompute τ_i 's and duration $\Delta T = \sum_1^N \tau_i$)

$\Delta t = \tau_i$

call EXPLICIT(Δt)

endfor

call OUTPUT if you want (if time $\geq t_{\text{out}}$)

Remarks: 1. $N=1, \nu=0$ is the explicit scheme itself, very convenient for debugging.

2. For fixed N , the smaller ν is the faster it runs (bigger ΔT , so fewer supersteps to get to t_{end})
but also less accurate.

So ν is a convenient gauge: close to $\nu=0$: more efficiency, less accuracy
further from $\nu=0$: more accuracy, higher cost

3. Experience shows that $N=5$ to 20 yields best results

4. Works on linear and nonlinear problems equally well

5. Works on 1, 2, 3 dimensions, higher efficiency in higher dimensions.

6. Beats the standard implicit schemes in efficiency and accuracy
with no extra programming!